## INFORMATION AND CODING THEORY Exercise Sheet 3

**Exercise 1. Data compression**. Determine whether the following codes are nonsingular, uniquely decodable or instantaneous:

- (a) The code {0,0}.
- (b) The code {0,010,01,10}.
- (c) The code  $\{10, 00, 11, 110\}$ .
- (d) The code {0, 10, 110, 111}.

**Exercise 2.** We consider a source  $\{x_i, p_i\}$  with  $i = 1, \dots, m$ . The symbols  $x_i$  (emitted with probabilities  $p_i$ ) are encoded in sequences using an alphabet of cardinality D, such that the decoding is instantaneous. For m = 6 and lengths of the codewords  $\{l_i\} = \{1, 1, 1, 2, 2, 3\}$ , find a lower bound for D. Is this code optimal?

**Exercise 3.** Huffman code. A source emits a random variable *X* which can take four values with probabilities  $(\frac{1}{10}, \frac{2}{10}, \frac{3}{10}, \frac{4}{10})$ .

- (a) Construct a binary Huffman code for *X*.
- (b) Construct a ternary Huffman code for *X*.
- (c) Construct a binary Shannon code for *X* and compare its expected length with the code of (a).

**Exercise 4.** Huffman code. A source emits a random variable *X* which can take four values with probabilities  $(\frac{1}{3}, \frac{1}{3}, \frac{1}{4}, \frac{1}{12})$ .

- (a) Construct a binary Huffman code for *X*.
- (b) Construct a binary Shannon code for *X* and compare it with the code of (a).

**Exercise 5.** We have a non-balanced coin with probability p to obtain "head" ("1") and probability 1 - p to obtain "tail" ("0"). Alice flips this coin as many times as needed to obtain "head" for the first time, and would like to communicate to Bob the number of flips k that were needed. A naive method is to send to Bob the sequence of the outcomes of the coin flips encoded in a chain of bits of length k, like  $000\cdots01$  (where 0 stands for "tail" and 1 for "head").

(a) What is the expected length of this naive code? Compare it with the entropy of the random variable *k*. In which case the naive code is optimal? Use

$$\sum_{i=0}^{\infty} a^i = \frac{1}{1-a}$$
$$\sum_{i=0}^{\infty} i a^i = \frac{a}{(1-a)^2}$$

(b) Alice decides now to encode her random variable *k* with a Shannon code, with the aim to approach H(k). Compare the expected lengths of the naive code and the Shannon code in the limit  $p \rightarrow 0$ .