INFORMATION, CODING, COMPUTING AND COMPLEXITY THEORY Exercise Sheet 1

**Exercise 1.** A discrete random variable *X* has a uniform probability distribution p(x) over the set  $\mathcal{X}$  of cardinality  $m = |\mathcal{X}|$ .

- (a) Give an example for  $\mathcal{X}$  and calculate p(x) for  $x \in \mathcal{X}$ .
- (b) Determine the entropy H(X) of X. If m = 64, what is H(X)?
- (c) How many bits are needed to enulerate the alphabet of  $\mathcal{X}$  without encoding, when m = 64?
- (d) How many symbols taken from a quaternary alphabet are necessary to enumerate X without encoding, when m = 64?
- (e) Show that H(X) is greater than the entropy of any other random variable *Y*, when *Y* takes values from the same set as  $\mathcal{X}$ . (*Read the reminder on Lagrange multipliers.*)

**Exercise 2.** A random variable *X* has an alphabet  $\mathcal{X} = \{A, B, C, D\}$ .

- (a) Calculate  $H(X_p)$  with associated probability distribution  $p(X) = \{\frac{1}{1}, \frac{1}{4}, \frac{1}{8}, \frac{1}{8}\}$ .
- (b) Calculate  $H(X_q)$  with associated probability distribution  $q(X) = \{\frac{1}{4}, \frac{1}{4}, \frac{1}{4}, \frac{1}{4}\}$ .
- (c) Consider the binary code defined by {C(X=A) = 0, C(X=B) = 10, C(X=C) = 110, C(X=D) = 111}.
  Calculate the expected length (in number of bits) of the codewords of *X*, when *X* obeys the distribution *p*(*X*) and *q*(*X*) respectively.
- (d) Compare the four results.

**Exercise 3.** A fair coin is flipped repeatedly until one obtains "head". Let the random variable *X* be given by the number of flips until one obtains "head" for the first time.

- (a) Calculate H(X).
- (b) Find a sequence of "yes/no" questions to determine the value of *X*. Compare the entropy of *X* with the expected length of the sequence of questions necessary to fully determine *X*.

Hint:

$$\sum_{n=1}^{\infty} r^n = \frac{r}{1-r},$$
$$\sum_{n=1}^{\infty} n \cdot r^n = \frac{r}{(1-r)^2}, r < 1.$$

**Exercise 4.** Let X and Y be two random variables. What is the inequality relating H(X) and H(Y) if,

- (a)  $Y = 2^X$ , with  $\mathcal{X} = \{0, 1, 2, 3, 4\}$  with associated probability distribution  $\{\frac{1}{2}, \frac{1}{4}, \frac{1}{8}, \frac{1}{16}, \frac{1}{16}\}$ .
- (b)  $Y = \cos(X)$ , with  $\mathcal{X} = \{0, \pi/2, \pi, 3\pi/2, 2\pi\}$  with associated probability distribution  $\{\frac{1}{3}, \frac{1}{3}, \frac{1}{12}, \frac{1}{12}, \frac{1}{6}\}$ .
- (c) Show that the entropy of any function f(X) of the random variable X is less than or equal to the entropy of X. In order to prove this, it is useful to calculate the joint entropy H(X, f(X)). In which case the is inequality  $H(f(X)) \le H(X)$  saturated?

**Exercise 5.** Let *X* and *Y* be two random variables, which take the values  $x_1, x_2, \dots, x_r$  and  $y_1, y_2, \dots, y_s$ . Furthermore, define a random variable Z = X + Y.

- (a) Show that H(Z|Y) = H(X|Y) and H(Z|X) = H(Y|X). Deduce that if *X* and *Y* are independently distributed, then  $H(X) \le H(Z)$  and  $H(Y) \le H(Z)$ . Thus, summing the two random variables can only increase the uncertainty.
- (b) Give an example for *X* and *Y* (which should be correlated) such that H(X) > H(Z) and H(Y) > H(Z).
- (c) In which case is the equality H(Z) = H(X) + H(Y) satisfied?

**Exercise 6. (Optional)** Give an example of a distribution of two random variables *X* and *Y* whose correlation coefficient *r* (as defined below) is zero, even though they are not independent. Show that in this case  $H(X:Y) \neq 0$ , which shows that the mutual entropy is a better measure of the dependence of *X* and *Y*.

Definition. Correlation coefficient

$$r = \frac{\langle x \cdot y \rangle - \langle x \rangle \cdot \langle y \rangle}{\sqrt{\langle x^2 \rangle - \langle x \rangle^2} \sqrt{\langle y^2 \rangle - \langle y \rangle^2}}$$

## Reminders

(a) How to change the base in logarithm:

$$\log_a b = \log_x b / \log_x a.$$

- (b) Lagrange theorem:
  - To maximize  $u(x_1, x_2, ..., x_n)$  with k constraints  $g_j(x_1, x_2, ..., x_n) = a_j, j = 1, \cdots, k$ .
  - Introduce the Lagrangian:  $L(x_1, ..., x_n, \lambda_1, ..., \lambda_k) = u(x_1, x_2, ..., x_n) + \sum_{j=1}^k \lambda_j [a_j - g_j(x_1, x_2, ..., x_n)].$
  - With the condition of an extremum being:  $\forall i, \frac{\partial L}{\partial x_i} = 0.$
  - If *u* is concave, the solution is a maximum.

The exercises and solutions are available at http://quic.ulb.ac.be/teaching